

# Stochastic Programming

## Lecture 7

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# Today's Class

- 1 Research Paper Presentations
- 2 Expected Value of Perfect Information (EVPI)
- 3 Value of the Stochastic Solution (VSS)
- 4 Basic Inequalities and EVPI vs. VSS
- 5 Relationship between VSS and EVPI
- 6 Bounds on EVPI and VSS

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# Research Paper Presentations

Starting in mid October we will begin reviewing research papers in the following tentative areas:

- Statistical sampling
- Multi-Stage Stochastic Linear Programming
- Bounding methods
- Integer L-shaped Method
- Probabilistic Constraints
- Robust Optimization

You can work in groups of 1 or 2. Please prioritize topics you are interested in and I will assign papers to cover.

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# Expected Value of Perfect Information (EVPI)

*Expected Value of Perfect Information (EVPI)* is a measure of how much value could be gained from predicting future outcomes.

## Definition:

$$\min z(x, \xi) = \min_x \{cx + \min_y \{q(\omega)^T y \mid Wy = h(\omega) - T(\omega)x, y \geq 0\} \mid Ax = b, x \geq 0\}$$

Assuming there is an optimal solution,  $\bar{x}(\xi)$ , for any  $\xi \in \Xi$  define two new problems

- *Wait and See Solution (WS)*:  $WS = E_{\xi}[\min z(x, \xi)] = E_{\xi}[z(\bar{x}(\xi), \xi)]$
- *2-Stage Recourse Problem (RP)*:  $RP = \min_x E_{\xi}z(x, \xi)$

# Expected Value of Perfect Information (EVPI)

The definition of EVPI is:

$$EVPI = RP - WS$$

Question: What is the EVPI for the newsvendor problem?

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# Value of the Stochastic Solution (VSS)

The *Value of the Stochastic Solution* (VSS) is the difference between the recourse problem (RP) and an “easy alternative”:

$$EV = \min_x \{z(x, E[\xi])\}$$

EV is the *mean value problem* (MV), with optimal solution  $\bar{x}$ .

$$EEV = E_{\xi}[z(\bar{x}, \xi)]$$

(EEV) is the expected value of using  $\bar{x}$  in the RP.

$$VSS = EEV - RP$$

# Value of the Stochastic Solution (VSS)

The *Value of the Stochastic Solution* (VSS) is:

$$VSS = EEV - RP$$

What is the VSS for the newsvendor problem?

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# Basic Inequalities and EVPI vs. VSS

Some useful propositions from Birge and Louveaux.....

Proposition 1 (Ch. 4 B & L):

$$WS \leq RP \leq EEV$$

Proposition 2 (Ch. 4 B & L):

*For stochastic programs with fixed objective coefficients, fixed  $T$ , and fixed  $W$ ,*

$$EV \leq WS$$

# Basic Inequalities and EVPI vs. VSS

## Proposition 3 (Ch. 4 B & L):

*Let  $x^*$  represent an optimal solution to the RP problem and let  $\bar{x}(\bar{\xi})$  be an optimal solution to the EV problem then*

$$RP \geq EEV + (x^* - \bar{x}(\xi))\eta$$

*where  $\eta \in \partial E_{\xi}[z(\bar{x}(\bar{\xi}, \xi))]$ , the subdifferential set of  $E_{\xi}[z(\bar{x}(\bar{\xi}, \xi))]$  at  $\bar{x}(\bar{\xi})$ .*

# Basic Inequalities and EVPI vs. VSS

Consider the following (slighted revised version of the recourse) problem:

Proposition 4 (Ch. 4 B & L):

$$\min z_u(x, \xi) = cx + \min\{qy \mid Wy \geq h(\xi) - Tx, y \geq 0\} \quad (1)$$

$$\text{s.t. } Ax = b \quad (2)$$

$$x \geq 0 \quad (3)$$

and the following definition of the RP

$$RP = \min_x \{E_\xi[z_u(x, \xi)]\}$$

Assume further that  $h(\omega)$  is bounded above by  $h_{\max}$ . Let  $x_{\max}$  be an optimal solution to  $z_u(x, h_{\max})$ . Then

$$RP \leq z_u(x_{\max}, h_{\max}).$$

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# Relationship between VSS and EVPI

EVPI and VSS measure different things and can have very different values. The following can be proved:

Proposition 5 (Ch. 4 B & L):

*For any stochastic program:*

$$EVPI \geq 0, \quad VSS \geq 0$$

*and for stochastic programs with fixed recourse and fixed objective function coefficients*

$$EVPI \leq EEV - EV, \quad VSS \leq EEV - EV$$



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# Bounds on EVPI and VSS

Consider a simplified version of a stochastic program where  $\xi = h(\omega)$  and  $\Xi$  is finite with scenarios  $\xi^k, k = 1, \dots, K$ . To develop the bounds we consider a *reference scenario*,  $\xi^r$ . Common reference scenarios are:

- expected value:  $\bar{\xi}$
- worst case scenario

Let  $p^r = P(\xi = \xi^r)$  and define the following to be the *PAIRS Subproblem*:

$$\min z^P(x, \xi^r, \xi^k) = cx + p^r qy(\xi^r) + (1 - p^r)qy(\xi^k) \quad (4)$$

$$\text{s.t.} \quad Ax = b \quad (5)$$

$$Wy(\xi^r) = \xi^r - Tx \quad (6)$$

$$Wy(\xi^k) = \xi^k - Tx \quad (7)$$

$$x, y \geq 0. \quad (8)$$

# Bounds on EVPI and VSS

Proposition 6 (Ch. 4 B & L):

*When the reference scenario is not in  $\Xi$ , then  $SPEV = WS$ , where  $SPEV$  is the sum of pairs expected values problem:*

$$SPEV = \frac{1}{1 - p^r} \sum_{k=1}^K p^k \min z^P(x, \xi^r, \xi^k)$$

# Bounds on EVPI and VSS

Proposition 7 (Ch. 4 B & L):

$$WS \leq SPEV \leq RP$$

Proposition 8 (Ch. 4 B & L):

$$RP \leq EPEV \leq EVRS$$

Proposition 9 (Ch. 4 B & L):

$$0 \leq EVRS - SPEV \leq VSS \leq EVRS - SPEV \leq EVRS - WS$$